A Vector Perturbation with User Selection for Multiuser MIMO Downlink

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Abstract—Recent works on multiuser MIMO study have shown that the linear growth of capacity in single user MIMO system can be translated to the multuser MIMO scenario as well. In this paper, we propose a method pursuing performance gain of vector perturbation in multiuser downlink systems. Instead of employing the maximum number of mobile users for communication, we use small part of them as virtual users for improving reliability of users participating communication. By controlling parameters of the virtual users including information and perturbation vector, we obtain considerable improvement in the effective SNR. Simulation results on the realistic multiuser MISO and MIMO downlink systems show that the proposed method brings substantial performance gain over the standard vector perturbation with marginal overhead in computations.

Index Terms—Multiuser downlink, vector perturbation, non-integer least square, sphere encoding, transmit precoding, virtual users.

I. INTRODUCTION

Recently, there has been growing interest in multiuser MIMO broadcast system where a basestation with multiple antennas transmits independent information to multiple mobile users, each with one or more receive antennas on the same frequency and time slots. Due to the fact that co-channel interference is hard to be managed by the receiver operation, and more importantly capacity of interference channel can be made equivalent to capacity without interference by judiciously controlling interference at the transmitter, pre-cancellation of the interference via the precoding at the transmitter has received much attention. It is now well known that dirty paper coding (DPC) [2] can achieve the capacity region of the multiuser MIMO downlink [3]–[5]. However, the shortcoming of the DPC is that system implementation leads to the high computational cost of successive encoding and decoding [6].

Lately, an approach referred to as vector perturbation technique [7] achieves the performance close to the capacity on multi-antenna multiuser downlink channel. By adding a deliberately designed perturbation vector into the information vector, the vector perturbation provides substantial improvement in the performance over the linear precoding methods such as channel inversion and regularized inversion with no virtual overhead in the receiver. Although the vector perturbation is a promising scheme for maximizing the sum rate of the MISO type broadcasting channel, we might need a strong protection against fading to meet the quality of service (QoS) of next-generation wireless systems.

In this paper, we put forth an approach improving the bit error rate performance by sacrificing the goal of maximizing the sum rate. To be specific, when the channel of certain users is in outage, we exploit these users to improve the reliability of the rest participating in communication. Towards this end, the proposed approach selects virtual users among users in outage to improve the signal-to interference-plus-noise ratio (SINR) of active users. It should be noted that, although the objective of proposed method and diversity technique [8] to resist the channel fading for improving reliability is the same, the means to achieve the goal is clearly different. While the diversity scheme accomplishes the goal by employing spatial or temporal replica of the transmitted signal, the proposed method attains the most out of the users whose contribution to the sum rate is trivial. Indeed, since parameters of sacrificing users including information data, perturbation vector, and even further the associated channel vector can be arbitrarily chosen, proper control of these parameters help improving the SINR of the mobile users participated in communication.

The rest of this paper is organized as follows. After a brief summary of system model and the vector perturbation technique in Section II, we present the proposed method in Section III. In section IV, we extend the proposed method to the multiuser MIMO downlink scenario. We discuss simulation results in Section V and provide conclusion in Section VI.

We briefly summarize notations used in this paper. We employ uppercase boldface letters for matrices and lowercase boldface for vectors. The superscripts $(\cdot)^H$ and $(\cdot)^T$ denote conjugate transpose and transpose, respectively. $\| \cdot \|$ indicates an $L_2$-norm of a vector. $\text{diag}(\cdot)$ is a diagonal matrix where nonzero elements exist only on the main diagonal of the matrix. $\text{CN}(m, \sigma^2)$ denotes a complex Gaussian random variable with mean $m$ and variance $\sigma^2$. $Q(x)$ denotes Q-function defined as $Q(x) = \int_{x}^{\infty} \frac{1}{\sqrt{2\pi}} \exp(-z^2/2) dz$.

II. SYSTEM MODEL AND PREVIOUS WORKS

In the multiuser downlink MISO system, a basestation composed of $M$ antennas transmits information to $K$ mobile users, each with a single receive antenna. Under the flat fading channel assumption, the received signal vector $y = [y_1, \cdots, y_K]^T$ can be described by

$$ y = Hx + n $$

where $H = [h_1^T, \cdots, h_K^T]^T$ is the $K \times M$ channel matrix where $h_i$ is the channel vector of the user $i$, $x =$
information becomes $\tilde{s} = s + \tau \ell$ where $\tau \in \mathbb{R}^+$ and $\ell = \ell_m + j\ell_m (\ell_m, \ell_m \in \mathbb{Z}^K)$ is a $K$-dimensional complex vector [7]. Similar to the regularized inversion technique, the transmitted signal of this approach is

$$x = \frac{P\tilde{s}}{\sqrt{\gamma}} = \frac{1}{\sqrt{\gamma}}H^H (HH^H + \alpha I)^{-1}s$$

(4)

where $\gamma = \|P\tilde{s}\|^2$. In this approach, $\ell$ minimizing $\gamma$ is chosen to boost the effective SNR and thus

$$\ell = \arg \min_{\ell} \|H^H (HH^H + \alpha I)^{-1}(s + \tau \ell')\|^2.$$ (5)

This integer least squares problem can be solved by a closest lattice point search, referred to as sphere encoder [7]. Note that since the vector perturbation searches $\ell$ from the set $\{ s + \tau \ell | \ell \in \mathbb{Z}^K + j\mathbb{Z}^K \}$ to generate minimum possible $\gamma$, $\gamma$ of the standard vector perturbation is always smaller than or equal to the regularized channel inversion where $\ell = 0$. Plugging (4) into (1), the received vector becomes $y = Hx + n \approx \frac{1}{\sqrt{\gamma}}s + n = \sqrt{\gamma}(s + \tau \ell) + n$. Since elements of $\ell$ are integers, $\tau \ell$ can be removed by the modulo operation [7] in the receiver and hence

$$y' = \frac{1}{\sqrt{\gamma}}(\gamma y \mod \tau) \approx \frac{s}{\sqrt{\gamma}} + n.$$ (6)

The salient feature of the vector perturbation lies on the use of integer based perturbation vector that can be easily removed by the modulo operation in the receiver, which provides considerable reduction of $\gamma$ (see Fig. 1).

III. VECTOR PERTURBATION WITH VIRTUAL USERS

A. Motivation

Although the vector perturbation method achieves bit error rate performance close to the sum capacity [4], in many practical scenarios where some mobiles are under deep fading over multiple codeblocks, performance loss for each user will be considerable. Recalling that the received vector of $j$th user is $y_j \approx \sqrt{\gamma} s_j + n_j$, the received SNR of user $j$ after the modulo operation becomes $\text{SNR}_j = \frac{E[|y_j|^2]}{E[|n_j|^2]}$. In fact, since $\gamma$ affects SNR of all users, whole mobile users will suffer performance degradation when the channel matrix is poorly conditioned due to the deep fading of certain users. An idea motivated by this observation is to exploit part of spatial degree of freedom as sacrificial lambs (virtual users) to enhance performance of the active users. Since the performance of virtual users is not our concern in this case, an aggressive control of their parameters is possible for maximizing performance gain of active users. In fact, the key distinctions of the proposed method over the standard vector perturbation are that 1) selection of virtual users when channels of certain users are in outage, 2) construction of well-conditioned $H$ matrix by using channel vectors of virtual users, and finally 3) selection of the non-integer perturbed input $s_j + \tau \ell_j$ for virtual users using a modified sphere encoding. For the analytic simplicity, we do not consider the modulo loss at the receiver in our sum rate
Thus, our result in this section might be slightly pessimistic.

**B. Virtual User Selection**

Our goal in this subsection is to select virtual users, only if at least one user is in outage, and then construct a channel matrix being used for the precoding. In the first stage, outage condition is tested for each user to select \(N_a\) active users \((N_a < K)\). Next, in case \(N_a \geq K\) (i.e., at least one virtual user is selected), \(K - N_a\) virtual channels are generated to replace the rows of sacrificing users with these.

To be specific, in the first stage, we check whether a user is in outage. For a given rate threshold \(\epsilon\), we say a user \(j\) is in \(\epsilon\)-outage if the rate \(R_j = \log (1 + \text{SNR}_j)\) of the user is smaller than \(\epsilon\). If no user is in \(\epsilon\)-outage, all users become the active users \((N_a = K)\). Otherwise, we choose at most \(K - N_a\) virtual users among users in outage.

Once at least one virtual user is chosen, our concern in \(H\) matrix construction is only on the rows \(h_k\) of virtual users \(k\). Instead of randomly generating \(h_k\), we select channels of mobile users not participating in real communication as candidates for virtual rows. That is, rows of sacrificing users in \(H\) are replaced by those of candidates. Since the generated channel matrix we use in the precoding is distinct from the original channel matrix \(H\), we henceforth denote it as \(\tilde{H}\).

Without loss of generality, we assume that users are sequenced in the following order: active users and virtual users. Then the top \(N_a\) rows of \(\tilde{H}\) corresponding to active users are equivalent to those of \(H\) and hence

\[
\tilde{H} \tilde{H}^H (\tilde{H} \tilde{H}^H)^{-1} = \begin{bmatrix} I_{N_a} & 0 \\ X_0 & X_1 \end{bmatrix}
\]

where \(X = [X_0 \ X_1]\) is the rows of virtual users for which we do not care. Clearly, there is no impact of virtual channel on the received signal of active users \(y_j\) for \(j = 1, \ldots, N_a\). After the precoding with \(\tilde{H}\), the received vector for active users is expressed as

\[
y_{a} = \frac{1}{\sqrt{n}} \tilde{s} + n.
\]

Let us now consider the virtual channel selection. Denoting \(\gamma\) of the proposed method as \(\gamma'\), \(\gamma' = ||\tilde{H}^H \left(\tilde{H} \tilde{H}^H\right)^{-1} \tilde{s}||^2\)

\[
= \tilde{s}^H \left(\tilde{H} \tilde{H}^H\right)^{-1} \tilde{s}.
\]

Also, using the eigenvalue decomposition of \(\tilde{H} \tilde{H}^H = Q \Lambda Q^H\), we have \(\gamma' = \tilde{s}^H Q \Lambda^{-1} Q^H \tilde{s}\). By letting \(\tilde{v} = Q^H \tilde{s}\), \(\gamma'\) is expressed as

\[
\gamma' = \tilde{v}^H \Lambda^{-1} \tilde{v} = \sum_{i=1}^{M} \frac{|\tilde{v}_i|^2}{\lambda_i}
\]

(9)

where \(\lambda_i\) is \(i\)th eigenvalue of \(\tilde{H}^H \tilde{H}\) and \(\tilde{v}_i = q_i^H \tilde{s}\) is \(i\)th element of \(\tilde{v}\) vector. \(q_i\) is the \(i\)th column of \(Q\). Since \(0 \leq \lambda_i \leq 8M\) [10] while \(\Sigma_{i=1}^{M} |\tilde{v}_i|^2 = ||\tilde{s}||^2\), variations among \(|\tilde{v}_i|^2\) is usually smaller than that of \(\lambda_i\) and thus the criterion to choose \(\tilde{H}\) among candidate matrix \(H_j\) can be approximated to

\[
\tilde{H} = \arg \min_{H_j} \sum_{i} \frac{1}{\lambda_i^{(j)}}
\]

(10)

where \(\tilde{H}_j\) is the \(j\)th candidate matrix including the rows \(h_j\) of virtual users \(j\) and \(\lambda_j^{(i)}\) is the \(i\)th eigenvalue of \(\tilde{H}_j \tilde{H}_j^H\). Note that in choosing \(\tilde{H}\) using (10), we need to find all the eigenvalues of \(\tilde{H}_j \tilde{H}_j^H\), which can be cumbersome for a large
dimensional matrix. Exploiting the approximation \( \sum_{i} \frac{1}{\lambda_i} = \frac{1}{\lambda_1} (1 + \frac{1}{\lambda_2} + \cdots + \frac{1}{\lambda_M}) \approx \frac{1}{\lambda_1} \), where \( \lambda_1 < \cdots < \lambda_M \), we obtain a simplified condition (max-min criterion) given by

\[
\hat{H} = \arg \max_{H} \lambda_{\gamma}^{(2)}.
\]  

(11)

By employing (11) instead of (10), computational overhead can be reduced substantially.²

Fig. 2(a) plots \( \gamma' \) as a function of SNR for the standard vector perturbation and proposed approaches employing (10) and (11). Although the average \( \gamma' \) using (10) is smaller than that obtained from (11) in low SNR regime, they become close when the SNR increases. This suggests that max-min criterion can provide cost-effective solution with negligible performance loss for mid and high SNR regime. In order to observe how \( \gamma' \) can provide cost-effective solution with negligible performance loss, we provide a simplified condition (max-min criterion) given by

\[
\gamma' = \min_{\ell_o \in \mathbb{U}^{M}} \min_{\ell_v \in \mathbb{C}^{N_v}} \| r - B_1 \ell_o - B_2 \ell_v \|_2^2
\]  

(16)

where \( B_1 \in \mathbb{C}^{M \times N_u} \) and \( B_2 \in \mathbb{C}^{M \times N_v} \) are the partitioned sub-matrices of \( B = [B_1, B_2] \).

The way to solve this problem consists of two steps. In the first step, optimal \( \ell_v \) for each \( \ell_o \) is computed using the standard least squares technique. Then \( \ell_v \) becomes a function of \( \ell_o \) (i.e., \( \ell_v = f(\ell_o) \)). Denoting \( g(\ell_o) = \min_{\ell_v \in \mathbb{C}^{N_v}} \| r - B_1 \ell_o - B_2 f(\ell_o) \|_2^2 \), (16) can be expressed as

\[
\gamma' = \min_{\ell_o \in \mathbb{U}^{M}} g(\ell_o).
\]  

(17)

In the second step, we solve this integer least squares problem (ILSP) using the sphere encoding. To be specific, for given \( \ell_o \) (i.e., when \( r - B_1 \ell_o \) is fixed), the optimal solution for the non-integer least squares problem becomes minimum norm solution [14], [15]

\[
\hat{\ell}_v = (B_2^H B_2)^{-1} B_2^H (r - B_1 \ell_o).
\]  

(18)

Plugging (18) into (16) and after some manipulation, we have

\[
\gamma' = \min_{\ell_o \in \mathbb{U}^{M}} \| (I - B_2 (B_2^H B_2)^{-1} B_2^H) r - (I - B_2 (B_2^H B_2)^{-1} B_2^H) B_1 \ell_o \|_2^2.
\]  

(19)

Let \( \Gamma = I - B_2 (B_2^H B_2)^{-1} B_2^H \), then (19) becomes

\[
\gamma' = \min_{\ell_o \in \mathbb{U}^{M}} \| \Gamma r - \Gamma B_1 \ell_o \|_2^2.
\]  

(20)

Note that \( \Gamma r \) is the projection of \( r \) onto the complement subspace of \( T = \text{span}(b_{2,i}) \) where \( b_{2,i} \) is the \( i \)th column vector of \( B_2 \). \( \Gamma B_1 \) can be interpreted in a similar way. Since we search \( \ell_v \) after projecting \( r \) and \( B_1 \ell_o \) vectors into the complement subspace of \( T \), the search space of the proposed method is reduced to \( N_v \)-dimensional integer lattice from the \( K \)-dimensional integer lattice (\( N_u < K \)). By letting \( r = \Gamma r \) and \( B_1 = \Gamma B_1 \) in (20), we have

\[
\gamma' = \min_{\ell_o \in \mathbb{U}^{M}} \| \Gamma r - \Gamma B_1 \ell_o \|_2^2.
\]  

(21)

Since this is a standard ILSP, \( \gamma' \) and \( \ell_v \) can be obtained by the sphere encoding. Once \( \ell_v \) is obtained, by plugging \( \ell_v \) back to (18), \( \ell_o \) is attained. Following simple example illustrates basic idea of the proposed approach.

**Example**: Suppose the channel of standard vector perturbation is given by

\[
H = \begin{bmatrix} 2 & 3 \\ 1 & 2 \end{bmatrix},
\]  

(22)

and the symbol vector chosen from 2-PAM is \( u = [1 \ -1]^T \), then \( \gamma' \) value obtained from (12) becomes 0.4396. Whereas, if

²Simple iteration algorithms to find the minimum eigenvalue of \( \hat{H} \hat{H}^H \) include inverse power method and shifted power method. Refer to [11], [12], [17] for detailed description of these.
a virtual user is introduced, then the second row of the virtual channel matrix can be designed to minimize \( \gamma' \) as

\[
\hat{H} = \begin{bmatrix} 2 & 3 \\ 2 & 1 \end{bmatrix},
\]

and the symbol vector chosen from 4-PAM (for ensuring 2 bits/channel use for both cases) is \( s' = [-\frac{1}{\sqrt{2}} 0]^T \). In this case, \( \gamma' \) value obtained from (20) becomes 0.3663, resulting in 20\% reduction of original value. When 2-PAM is employed, \( \gamma' \) value becomes 0.2731, but the information rate is reduced to half.

To analyze benefit of the proposed method over the standard vector perturbation, we consider the error probability using the QPSK modulation. Denoting the error probability of the proposed method and standard vector perturbation as \( P' \) and \( P_e \), respectively, then the relationship between two can be described as \( P' = \beta P_e \) where \( \beta \) is the attenuation factor indicating the improvement of the error performance (the decibel gain in performance is \(-10 \log(\beta)\)).

**Lemma 1:** The attenuation factor \( \beta \) is lower bounded by

\[
\beta > \sqrt{\frac{2}{\pi^2 \gamma c}} \left[ 1 - \frac{\gamma'}{\gamma} \exp \left( -\frac{c}{2\gamma} \left( 1 - \frac{\gamma'}{\gamma} \right) \right) \right] \tag{24}
\]

where \( c = \frac{\gamma}{\sqrt{2}} \).

**Proof:** Recalling that the error probability of QPSK modulation is \( Q(\sqrt{2SNR}) \), \( P_e \) and \( P'_e \) become \( Q \left( \sqrt{\frac{c}{\gamma}} \right) \) and \( Q \left( \sqrt{\frac{c}{\gamma'}} \right) \), where \( c = \frac{\gamma}{\sqrt{2}} \). Using the properties of Q-function [16], we have

\[
P_e \leq \frac{1}{2} \exp \left( -\frac{c}{2\gamma} \right) \tag{25}
\]

\[
P'_e > \frac{1}{\sqrt{2\pi c}} \frac{\gamma'}{\gamma} \left[ 1 - \frac{\gamma'}{\gamma} \right] \exp \left( -\frac{c}{2\gamma} \right). \tag{26}
\]

Using (25) and (26) together with \( P' = \beta P_e \), we get (24). □

This result states that the gain of the proposed method strongly depends on \( \gamma' \). The overall operation of the proposed method is summarized in Table I.

### D. Lower Bound of the Sum Rate Loss

The sum rate achieved by the standard vector perturbation is

\[
R_{org} = \sum_{j=1}^{K} R_j = \sum_{j=1}^{K} \log(1 + SNR_j) \tag{27}
\]

where \( SNR_j = \frac{E||s||^2}{E||n||^2} \). Without loss of generality, we assume that \( R_1 > \cdots > R_K \). Then the virtual user selection strategy searches at most \( K - N_{a,min} \) users in outage (\( R_i < \epsilon \)).

**Lemma 2:** For given \( \epsilon > 0 \), the sum rate of the proposed method satisfies

\[
R_{prop} > R_{org} = (K - N_{a,min})\epsilon. \tag{28}
\]

**Proof:** It is clear from (27) that \( R_{org} = \sum_{j=1}^{N_{a,min}} R_j + \sum_{j=N_{a,min}}^{K} R_j \). Also, the sum rate of the proposed method is

\[
R_{prop} = \sum_{j=1}^{N_{a,min}} \log(1 + SNR'_j) \text{ where } SNR'_j = \frac{E||s'||^2}{E||n||^2}; \tag{29}
\]

Since \( N_a \geq N_{a,min} \) and \( \gamma > \gamma' \), we have

\[
R_{prop} \geq \sum_{j=1}^{N_{a,min}} R_j = R_{org} - \sum_{j=N_{a,min}+1}^{K} R_j
\]

\[
> R_{org} - (K - N_{a,min}) \epsilon.
\]

In particular, if \( N_{a,min} = K - 1 \), then \( R_{prop} > R_{org} - \epsilon \). □

Note that since (28) holds for any channel realization, the lemma will also hold in ergodic sense (average over all possible channel realizations). Note also that due to the reduction of \( \gamma \) and sporadic use of virtual users, the sum rate loss will be smaller than \((K - N_{a,min})\epsilon\) in practice.

### E. Comments on Complexity

In this subsection, we compare the complexity of the proposed method and the standard vector perturbation. Note that since the modulo operation at the receiver is common, we focus only on computations associated with the transmitter precoding. While the major operation of the conventional approach is the sphere encoding to compute perturbation vector \( \ell \), user selection step and modified sphere encoding
are additionally required for the proposed method. Denoting the complexity associated with user selection, $\ell_v$ computation in (18), and the sphere encoding as $C_{\text{SE}}$, respectively, then the number of required floating-point operations (flops) for each step is as follows [17]:

- $C_{\ell_v}$ requires $\frac{4}{3}M^3$ flops for inverting of $\tilde{H}_j \tilde{H}_j^H$ matrix, $k = 2M^2 + 2M^2 + 4M$ flops for computing the minimum eigenvalue using an inverse power method ($k$ is the iteration number). These operations need to be repeated for $\left( \frac{N_v}{N_c} \right)$ times for choosing $N_v$ virtual users among $N_c$ candidates.
- $C_{\ell_v}$ requires $\frac{4}{3}N_v^3 + 2MN_v^2$ for computing $(B_v^H B_v)^{-1}$ and $2MN_v^2 + 2MN_v + 2MN_a + M$ for the matrix multiplication in (18).

In contrast to the operations we just described, the complexity associated with the perturbation vector optimization is hard to be quantified mainly because the computational complexity of the sphere encoding is non-deterministic. As a popular way to quantify the complexity, the average number of nodes visited has been considered [18], [19]. The lower bound on the complexity of the sphere encoding $C_{\text{SE}}$ is [18]

$$C_{\text{SE}} \geq \frac{\beta \psi N_v - 1}{\sqrt{\beta - 1}}$$  \hspace{1cm} (30)

where $\beta$ is a modulation order, $N_v$ is the number of search dimension, and $\nu$ is the complexity component given by $\nu = \frac{1}{2} \left( 1 + \frac{2(\beta-1)}{\beta \omega} \right)$ where $\omega$ is a constant. Noting that $C_{\text{SE}}$ increases exponentially with the search dimension $N_v$ and $N_a$ of the modified sphere encoding is smaller than that of the standard sphere encoding ($K \geq N_a$), we can argue that the complexity associated with the modified sphere encoding is smaller than that associated with the sphere encoding of the standard vector perturbation. However, since the proposed method requires two sphere encoding operations (standard sphere encoding in (12) and modified sphere encoding in (21)), the overall complexity (sum of two sphere encoding) of the proposed method is larger than the standard vector perturbation. Note that due to the variation of complexity and latency, hardware implementation of the sphere encoding might be a difficult task. As a way to get around the latency and complexity issue, the fixed-complexity sphere encoding (FSE) might be considered [20], [21]. This together with the fact that extra operations are required only when virtual users are selected due to the $\epsilon$-outage, we expect that the additional computations of the proposed method are fairly moderate. Table II summarizes flops of the proposed method and standard vector perturbation. We observe that the additional complexity of the proposed method over the standard vector perturbation are 22\% ($N_v = 1$) and 26\% ($N_v = 2$), respectively. We also observe that the complexity of the sphere encoding with $N_v = 2$ is smaller than that with $N_v = 1$ since the search dimension is reduced from 3 to 2. However, due to the increase in complexity of user selection part, overall complexity for $N_v = 1$ and 2 is more or less similar.

### F. Scheduling for fairness

An important issue to be considered, when user channels are heterogeneous (i.e, there are large variations in SNR among users) or some users are under stringent delay constraint, is to guarantee fairness among users. Note that if the user selection threshold is set to $\epsilon = \frac{1}{M} R_{\text{avg}}$, all symbols of worst users are sacrificed for the performance improvement of active users so that no fairness is guaranteed for the sacrifice user at all. Whereas, if $\epsilon = 0$, the proposed method returns to the standard vector perturbation so that strict fairness is ensured among users via a simple round robin scheduling. As an approach making the compromise between the fairness and the sum rate, proportional fair scheduling (PFS) has been popularly considered in practice [22]. In each scheduling instant, the PFS computes the ratio of the instantaneous to the average rate for each user and then selects users associated with the maximal ratio ($k^*(t) = \arg \max_k \frac{R_k}{\bar{R}}$, $\bar{R}$). Even in this case, since the supportable rate $R_k$ for the PFS decision is computed using channel vectors $H_k$, only, we can combine the proposed method and the PFS to improve performance of active users. As mentioned, since all parameters of sacrificing users (information data, perturbation vector, and virtual channel vector) can be arbitrarily chosen to improve the SINR of the system, considerable performance gain can be obtained with negligible impact on the fairness.

We summarize the modified PFS algorithm in Table III. Due to the fact that the virtual user decision is performed after the PFS scheduling, the weight $\mu_k$ of the sacrificing user will increase, thereby providing better chance for this user in the next scheduling instance.$^3$

### IV. EXTENSION TO MULTIUSER MIMO DOWNLINK SCENARIO

So far, we have considered multiuser MISO scenario where the worst user under deep fading is sacrificed to improve the performance of other users. We can readily extend the proposed method to the multiuser MIMO downlink scenario. Note that, since each user has multiple streams in this scenario, we do sacrifice worst stream instead of user. In the multiuser MIMO downlink system, a basestation equipped with $N_t$ antennas transmits information to $K$ users. Each user has $N_{r,j}$ receive antennas and thus the total number of receive antennas is $N_r = \sum_{j=1}^{K} N_{r,j}$. In the sequel, we will use the notation $\{N_{r,1}, \ldots, N_{r,K}\} \times N_t$ for this multiuser MIMO

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$^3$For ensuring the strict sense fairness, one might add the condition for the minimal service requirement.
becomes the right choice for the precoding matrix $T$ at the system. Denoting $\tilde{T}$
where $\tilde{\nu}_j$ is the optimization function of the $\tilde{\nu}_j$ to the $\tilde{\nu}_j$.

The main idea of the BD is to employ the beamforming MIMO channel into the single user MIMO channel, the transmit power suppression caused by the zero forcing operation of the BD.

As a well-known approach for converting the multiuser MIMO channel into the single user MIMO channel, the SVD of $H$ can be applied to compensate for the transmit power allocation from the transmitter to the $j$th user, $x_j$ is the $N_{r,j} \times 1$ transmit signal vector of user $j$, and $n_j$ is the normalized complex Gaussian noise vector with entries distributed according to $CN(0,1)$. Clearly, $T_j$ should be a nullspace of $\tilde{H}_j = [H_1^T \cdots H_{j-1}^T \cdots H_{j+1}^T \cdots H_{K}^T]^T$. If we denote the SVD of $\tilde{H}_j$ as $\tilde{H}_j = \tilde{U}_j \tilde{\Sigma}_j \tilde{V}_j^{(0)}$, then the right singular vectors $\tilde{V}_j^{(0)}$ form an orthogonal basis for the null space of $\tilde{H}_j$. Thus, $\tilde{V}_j^{(0)}$ becomes the right choice for the precoding matrix $T_j$ and

$$y_j = H_j^* \sum_{j=1}^{K} T_j x_j + n_j \quad (31)$$

where $H_j$ is the $N_{r,j} \times N_t$ channel matrix from the transmitter to the $j$th user, $x_j$ is the $N_{r,j} \times 1$ transmit signal vector of user $j$, and $n_j$ is the normalized complex Gaussian noise vector with entries distributed according to $CN(0,1)$.

As a well-known approach for converting the multiuser MIMO channel into the single user MIMO channel, the block diagonalization (BD) algorithm [23] can be employed. The main idea of the BD is to employ the beamforming vectors annihilating all multiuser interference. The preceding constraint to remove multiuser interference is

$$H_i T_j = 0 \quad \text{for } i = 1, \cdots, j - 1, j + 1, \cdots, K. \quad (32)$$

Clearly, $T_j$ should be a nullspace of $\tilde{H}_j = [H_1^T \cdots H_{j-1}^T \cdots H_{j+1}^T \cdots H_{K}^T]^T$. If we denote the SVD of $\tilde{H}_j$ as $\tilde{H}_j = \tilde{U}_j \tilde{\Sigma}_j \tilde{V}_j^{(0)}$, then the right singular vectors $\tilde{V}_j^{(0)}$ form an orthogonal basis for the null space of $\tilde{H}_j$. Thus, $\tilde{V}_j^{(0)}$ becomes the right choice for the precoding matrix $T_j$ and

$$y_j = H_{eff,j} x_j + n_j = H_j T_j x_j + n_j = H_j \tilde{V}_j^{(0)} x_j + n_j \quad (33)$$

where $H_{eff,j}$ is the effective channel matrix for the $j$th user. By precoding with $\tilde{V}_j^{(0)}$, the multiuser interference is completely removed and thus the $j$th user experiences a point-to-point MIMO link [23].

After the BD preprocessing, the vector perturbation technique can be applied to compensate for the transmit power suppression caused by the zero forcing operation of the BD. Since it is straightforward to apply the vector perturbation technique for the system in (33), we skip the details. The optimization function of the $j$th user is

$$\ell_j = \arg \min_{\ell_j} \left\| H_{eff,j}^{H}(s_j + \tau \ell_j) \right\|^2 \quad (34)$$

where $Z$ is the integer set and $N_{r,j}$ is the number of transmit streams for each user. The achievable rate of this MU-MIMO scheme is [24], [25]

$$R_{\text{org-MIMO}} = \sum_{j=1}^{K} R_{j,\text{org}} = \sum_{j=1}^{K} \sum_{k=1}^{N_{r,j}} \log_2 (1 + \text{SNR}_{j,k})$$

$$= \sum_{j=1}^{K} \sum_{k=1}^{N_{r,j}} \left( 1 + \frac{\rho \lambda_k}{N_{r,j}} \right) \quad (35)$$

where $\rho = \frac{P_T}{\sigma_k^2}$ and $\lambda_k (k = 1, \cdots, N_{r,j})$ is the singular value of $H_{eff,j}$. Same as the MISO scenario, we can employ the proposed scheme for the selection of virtual streams when effective channels of certain streams are in outage. Since the system is decomposed into $K$-independent single user MIMO systems depending on the QoS of each user, we might use separate outage threshold $\epsilon_j$ for each user’s effective channel matrix $H_{eff,j}$. The achievable sum rate for the proposed MU-MIMO scheme is given by (see Appendix A)

$$R_{\text{prop-MIMO}} = \sum_{j=1}^{K} \sum_{l=1}^{N_a} \log \left( 1 + \frac{\rho \xi_j^2}{\min_{l_m} \left( \sum_{m=1}^{N_{r,j}} \rho \xi_j^2 \right)} \right) \quad (36)$$
the figure, the sum rate difference between two approaches is small and roughly constant. Therefore, as the SNR increases, percent loss ($\frac{R_{\text{org}} - R_{\text{prop}}}{R_{\text{org}}}$) of the sum rate decreases, resulting in about 12% loss at 20 dB SNR. In Fig. 3(b), we compare the BER performance of the proposed method along with standard vector perturbation and others. The performance gain, achieved at the expense of the sum rate loss, is noticeable so that the gain at $10^{-2}$ BER is around 2.5 dB. We also observe that the lower bound obtained by (24) is tight for mid and high SNR regime so that we can easily predict the achievable performance gain in the design stage. Overall, we observe that the proposed method with $\alpha = 0.1$ achieves higher sum rate than that with $\alpha = 0.15$, resulting in 7% gain at 20 dB SNR at the expense of 0.5 dB loss in performance.

In Fig. 4, we plot the sum rate and BER performance of the proposed schemes for $N_v = 1$ and 2. As expected, the performance of the proposed method with $N_v = 2$ outperforms that with $N_v = 1$ and standard vector perturbation, resulting in about 1.5 dB and 3 dB gain at $\text{BER} = 10^{-3}$, respectively. Since the performance trades off the sum rate, $N_v = 2$ scenario incurs additional 15% loss in the sum rate.

So far, we have assumed that the base station has knowledge of full channel state information (CSI). In Fig. 5, we investigate the performance of the proposed method along with the
standard vector perturbation in the presence of channel estimation error. Since the mismatch between the actual CSI and the estimated CSI is unavoidable in a real communication, and this might result in degradation in performance, it is of importance to investigate the effect of channel estimation error. In our simulation, we use an additive channel estimation error model where $H = H_{est} + H_{err}$ where $H$, $H_{est}$ and $H_{err}$ represent the true channel matrix, the estimated channel matrix and the estimated error matrix, respectively. We assume that $H_{err}$ is uncorrelated with $H_{est}$ and $s$, and $H_{err}$ has i.i.d. elements with zero mean and the estimation error variance $\sigma^2_{e,h}$. In this setup, it is clear that the estimation error becomes dominant factor limiting performance improvement in high SNR. Although both schemes show error floors when $\sigma^2_{e,h} = 0.05$, we observe that the proposed scheme is more robust to the estimation errors than the standard vector perturbation. For example, the performance gain at $10^{-3}$ BER of the proposed scheme is about 3 dB over the standard vector perturbation when $\sigma^2_{e,h} = 0.01$, while the gain is around 2 dB for $\sigma^2_{e,h} = 0$.

Finally, we consider the $\{4,4\} \times 8$ MU-MIMO system with the QPSK modulation. As mentioned, the multiuser system is converted into the single user MIMO system by the BD operation so that we might need an individual outage thresholds $\epsilon_1$ and $\epsilon_2$. For simplicity, we set $\epsilon_1 = \epsilon_2 = 0.15R_{1_{org}}$ in our simulation and hence the lower bound of the sum rate becomes $\sum_{j=1}^{2} (R_{j_{org}} - (N_{r,j} - N_{a_{min}})\epsilon_1)$. We observe form the Fig. 6(a) that the sum rate difference between two approaches is moderate and further decreases as SNR increases. Whereas, as shown in Fig. 6(b), the proposed method brings considerable gain over the BD and the BD with vector perturbation, resulting in 8 dB and 2.5 dB gain at $10^{-2}$ BER.

VI. CONCLUSION

In this work, we investigated an approach achieving robustness of multiuser downlink in the vector perturbation. Motivated by the fact that whole mobile users in a link suffer performance degradation since $\gamma$ of the vector perturbation affects SNR of all users alike, we sacrificed small part of users in outage to improve the bit error performance of the rest. This can eventually strengthen the ability to meet the QoS of the practical mobile communication system. The benefit of the proposed method, on top of the enhanced reliability, lies in the point that the modification of transmit precoding does not affect the receiver operation at all and the additional complexity of the proposed method at the transmitter is fairly moderate. Future research direction may include the complexity reduction of the transmission precoding and application of
the proposed method into the codebook based systems.

APPENDIX A

ACHIEVABLE RATE ANALYSIS OF THE PROPOSED MU-MIMO SCHEME

Recall that the received signal $y_j$ is given by $y_j = H_{\text{eff},j}x + n_j$ where $x_j = \frac{1}{\sqrt{N_j}}H_{\text{eff},j}\tilde{s}_j$. Define $H_{\text{eff},j} = \sum_{l=1}^{N_j} \lambda_l \mathbf{V}_l^H$, where $U_l = [u_{1,l} \cdots u_{N_j,l}]$, $V_j = [v_1 \cdots v_{N_j}]$, and $\Lambda = \text{diag}(\lambda_1 \cdots \lambda_{N_j})$. The normalized scalar factor $\gamma'_j$ is given by

$$\gamma'_j = \|H_{\text{eff},j}^{-1}s_j\|^2 = s_j^H\left(H_{\text{eff},j}H_{\text{eff},j}^H\right)^{-1}s_j$$

$$= s_j^H\Lambda^{-2}U_j^H\Lambda^{-2}U_j s_j = \sum_{l=1}^{N_j} \frac{1}{\mu_l} \xi_l,$$

(37)

where $\mu_l = \frac{1}{\xi_l} > 0$ and $\xi_l = \|u_l^H s_j\|^2$.

The received signal-to-noise-ratio (SNR) of each stream $\xi_l$ can be represented by $\text{SNR}_l = \frac{\rho \xi_l^2}{\sum_{m=1}^{N_j} \mu_m^2 \xi_m^2}$, where $\rho = P_a/\sigma_n^2$. Thus, the achievable rate of the $j$th user $R_j$ is given by

$$R_j = \sum_{l=1}^{N_j} \log(1 + \text{SNR}_l) = \sum_{l=1}^{N_j} \log \left(1 + \frac{\rho \xi_l^2}{\sum_{m=1}^{N_j} \mu_m^2 \xi_m^2} \right).$$

Note that the vector precoding is not applied in (38). The effect of vector perturbation is to force $\xi_l$ to minimize $\gamma'_j$ so that $\tilde{s}_j$ can be (coarsely) oriented in the coordinate system defined by $u_1, \cdots, u_{N_j}$ [7].

By searching the proper precoding vector and controlling $\xi_m$ to minimize $\gamma'_j$, we obtain an upper bound on achievable rate of the proposed scheme for each user as

$$R_{j,\text{prop}} = \sum_{l=1}^{N_j} \log \left(1 + \frac{\rho \xi_l^2}{\min_{\xi_m} \sum_{m=1}^{N_j} \mu_m^2 \xi_m^2} \right).$$

(39)

From (39), we need to solve for $\xi_l = \arg\min_{\xi_m} \sum_{m=1}^{N_j} \mu_m^2 \xi_m^2$.

By the arithmetic-geometric mean inequality, solution occurs when $\mu_l^2 \xi_l^2 = \cdots = \mu_{N_j}^2 \xi_{N_j}^2 = \omega_0^2$ for an arbitrary constant $\omega_0^2 > 0$. Note that $\xi_l = \|u_l^H \tilde{s}_j\|^2$ cannot be zero since $\tau$ is chosen large enough. Therefore, $R_{j,\text{prop}}$ can be represented by

$$R_{j,\text{prop}} = \sum_{l=1}^{N_j} \log \left(1 + \frac{\rho \omega_0^2}{\mu_l^2 N_j \omega_0^2} \right) = \sum_{l=1}^{N_j} \log \left(1 + \frac{\rho}{\mu_l^2 N_j} \right).$$

(40)

We obtain (40) by substituting $\mu_l$ for $\frac{1}{\xi_l}$. In summary, the upper bound of achievable sum rate for the proposed scheme becomes

$$R_{\text{prop-MIMO}} = \sum_{j=1}^{K} R_{j,\text{prop}}.$$

(41)
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